

Learning Biological Regulatory Networks from Time Series with LFIT and Application to Phytoplankton

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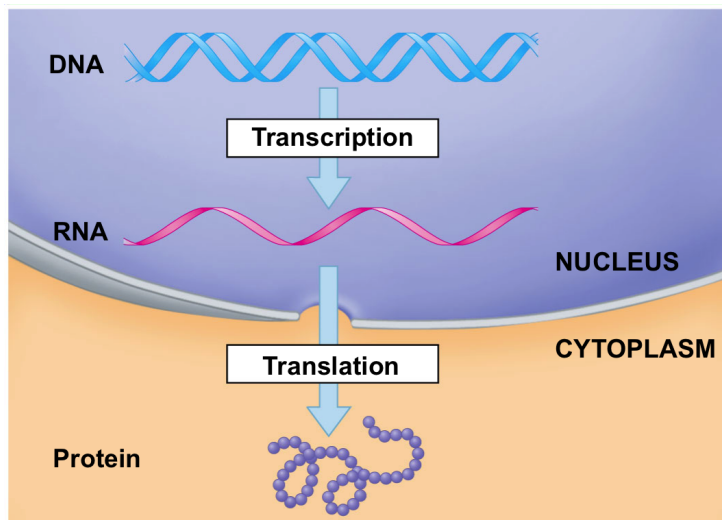
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MIAT Seminar

Joint work with: Tony RIBEIRO (Centrale Nantes), Morgan MAGNIN (Centrale Nantes), Katsumi INOUE (NII, Japan), Madeleine EYRAUD (CNRS), Sébastien LEFEBVRE (Univ. Lille), Cédric LHOSSAINE (Univ. Lille)

Outline

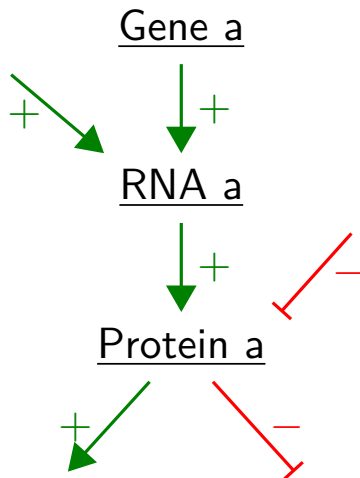
- Biological regulatory networks
- LFIT: an approach to learn a discrete model from its stage graph
- Heuristic for noisy/incomplete data
- Application to phytoplankton monitoring

Preliminary Abstraction

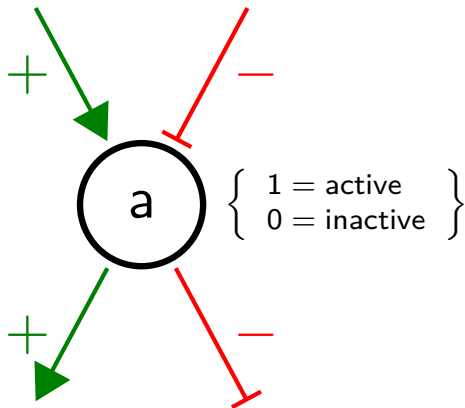


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Preliminary Abstraction



Preliminary Abstraction



Biological Regulatory Networks

Discrete Networks / Thomas Modeling

[Kauffman, *Journal of Theoretical Biology*, 1969]

[Thomas, *Journal of Theoretical Biology*, 1973]

- A set of components $N = \{a, b, z\}$

a

z

b

Discrete Networks / Thomas Modeling

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- A set of components $N = \{a, b, z\}$
- A discrete domain for each component $\text{dom}(a) = \{0, 1, 2\}$

$\{0, 1, 2\}$

a

z

$\{0, 1\}$

b

$\{0, 1\}$

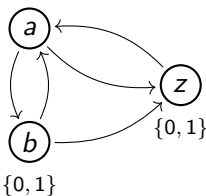
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- A set of components $N = \{a, b, z\}$
- A discrete domain for each component $\text{dom}(a) = \{0, 1, 2\}$
- Discrete parameters / evolution functions $f_a : \mathcal{S} \rightarrow \text{dom}(a)$

$\{0, 1, 2\}$



a	f_b
0	0
1	1
2	1

z	b	f_a
0	0	1
0	1	0
1	0	1
1	1	2

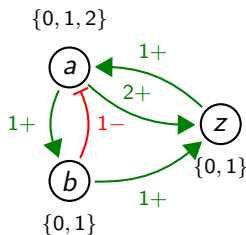
a	b	f_z
0	0	0
0	1	0
1	0	0
1	1	0
2	0	0
2	1	1

Discrete Networks / Thomas Modeling

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- A set of components $N = \{a, b, z\}$
- A discrete domain for each component $\text{dom}(a) = \{0, 1, 2\}$
- Discrete parameters / evolution functions $f_a : \mathcal{S} \rightarrow \text{dom}(a)$
- Signs & thresholds on the edges (redundant) $a \xrightarrow{2+} z$



a	f_b	z	b	f_a	a	b	f_z
0	0	0	0	1	0	0	0
1	1	0	1	0	0	1	0
2	1	1	0	1	1	0	0
		1	1	2	1	1	0
					2	0	0
					2	1	1

State Graph

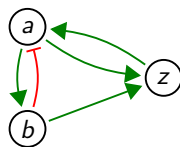
The state graph depicts explicitly the whole dynamics

abz

000 010 001 011

100 110 101 111

200 210 201 211



+ f_a, f_b, f_c

State Graph

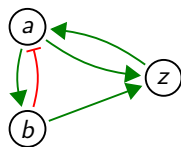
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abz

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100 $\longrightarrow$ 110 101 111

200 210 201 211

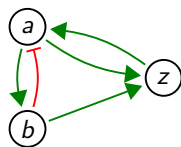
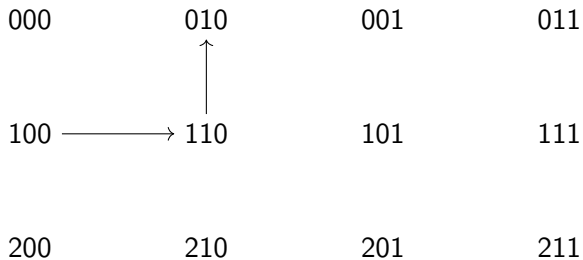


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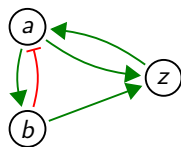
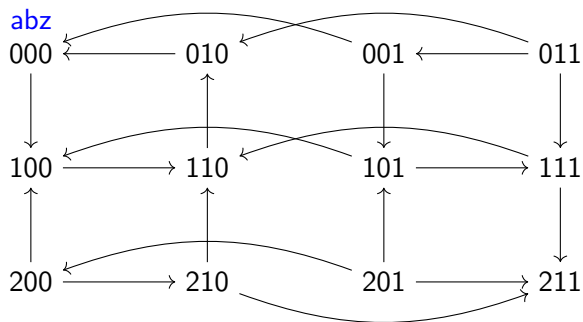
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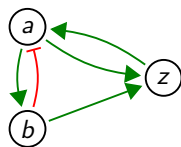
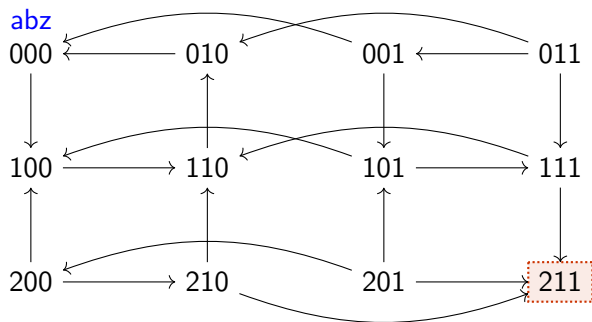
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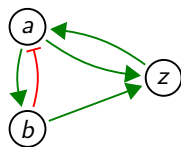
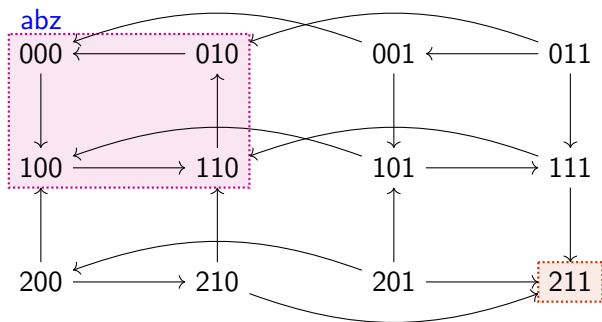


+ f_a, f_b, f_c

- **Stable state** = state with no successors

State Graph

The state graph depicts explicitly the whole dynamics

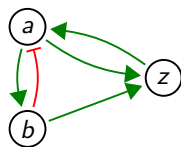
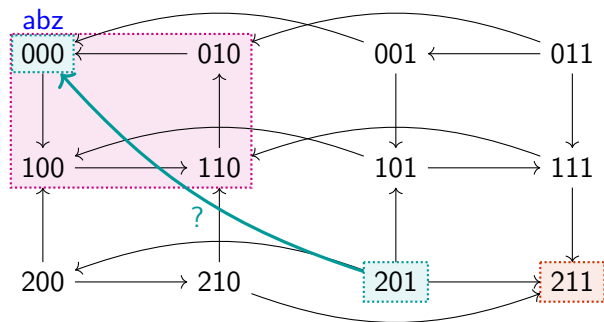


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- **Stable state** = state with no successors
- **Complex attractor** = minimal loop or composition of loops from which the dynamics cannot escape

State Graph

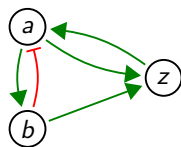
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- **Reachability** = from **201**, can I reach **000**?

The state graph depicts explicitly the whole dynamics

 $+ f_a, f_b, f_c$

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Semantics



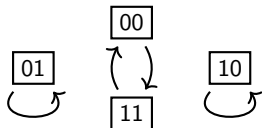
$$f_a = \neg b$$

$$f_b = \neg a$$

b	f_a
0	1
1	0

a	f_b
0	1
1	0

State transitions differ according to the update semantics used:



Synchronous

- **Synchronous:** all variables are updated

Semantics



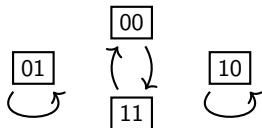
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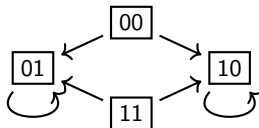
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State transitions differ according to the update semantics used:



Synchronous



Asynchronous

- **Synchronous:** all variables are updated
- **Asynchronous:** only one variable is updated

Semantics



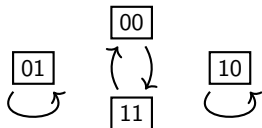
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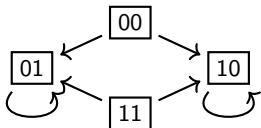
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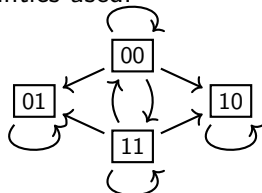
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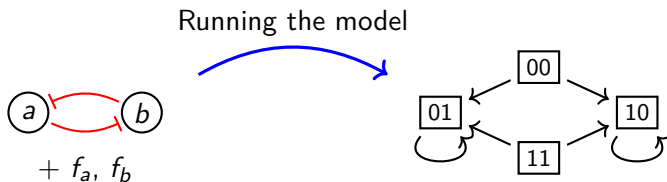
Asynchronous



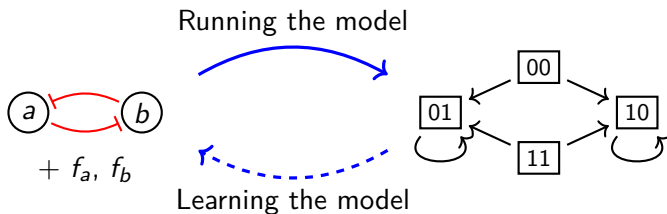
General

- **Synchronous**: all variables are updated
- **Asynchronous**: only one variable is updated
- **General**: any number of variables can be updated

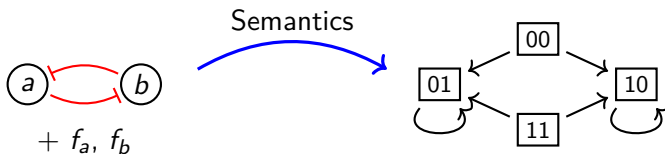
Learning from the State Graph



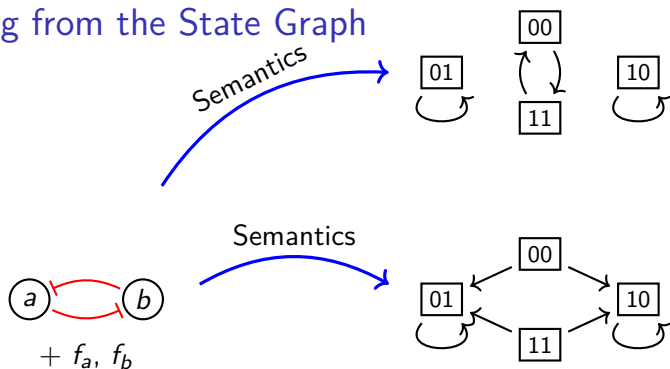
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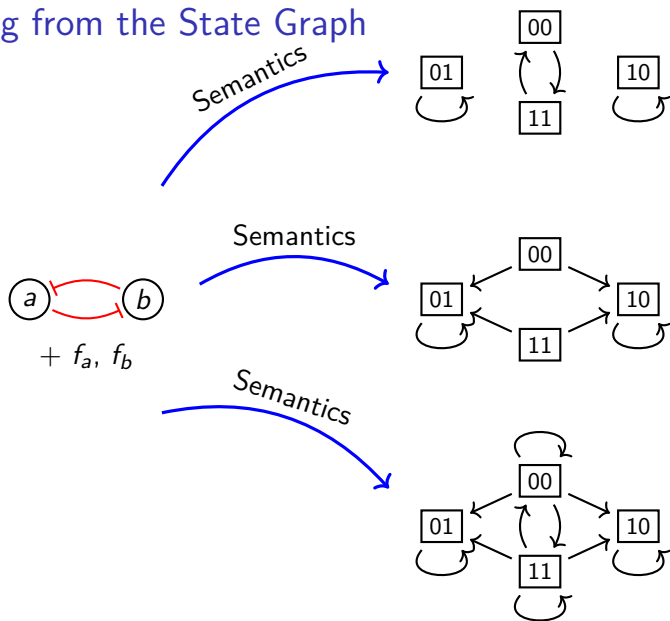
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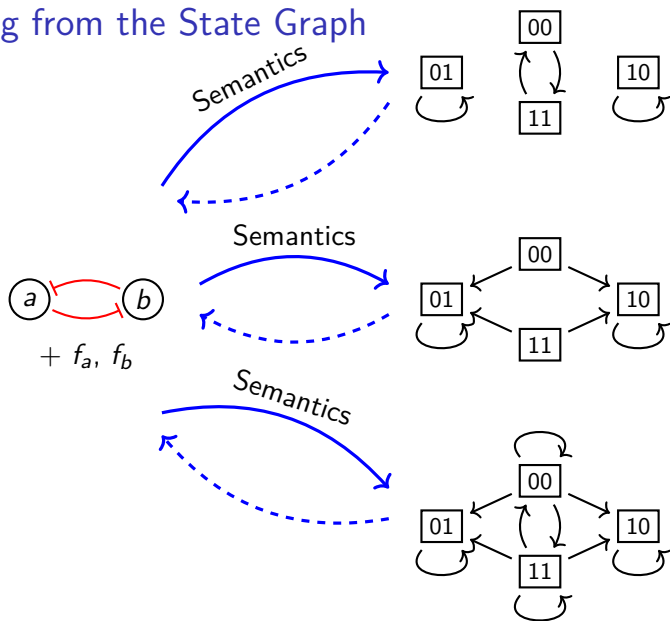
Learning from the State Graph



Learning from the State Graph



Learning from the State Graph



Logic Rules

A logic program is a set of logic rules.

It is an alternative representation of biological networks.

$$a_1 \leftarrow a_0, b_0, c_2.$$

If a and b are at level 0 and c is at level 2, then a can change its value to 1.

$$a_1 \leftarrow c_2.$$

Whenever c is at level 2, a can change its value to 1.

$$a_1 \leftarrow .$$

a can change its value to 1 anytime.

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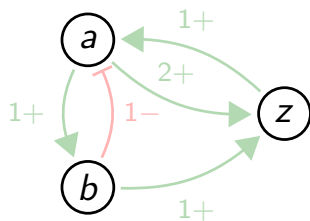
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The same notion of **semantics** applies.

Discrete Models as Logic Programs

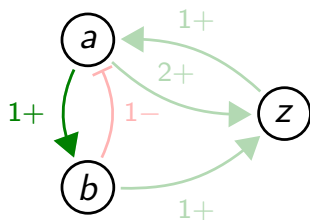
Discrete model:



Logic program:

Discrete Models as Logic Programs

Discrete model:



Logic program:

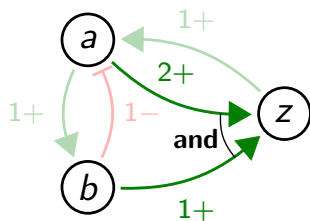
$$b_1 \leftarrow a_1.$$

$$b_1 \leftarrow a_2.$$

$$b_0 \leftarrow a_0.$$

Discrete Models as Logic Programs

Discrete model:



Logic program:

$$b_1 \leftarrow a_1.$$

$$b_1 \leftarrow a_2.$$

$$b_0 \leftarrow a_0.$$

$$z_1 \leftarrow a_2 \wedge b_1.$$

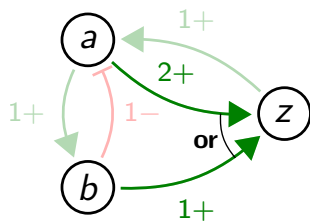
$$z_0 \leftarrow a_0.$$

$$z_0 \leftarrow a_1.$$

$$z_0 \leftarrow b_0.$$

Discrete Models as Logic Programs

Discrete model:



Logic program:

$$b_1 \leftarrow a_1.$$

$$b_1 \leftarrow a_2.$$

$$b_0 \leftarrow a_0.$$

$$z_1 \leftarrow a_2.$$

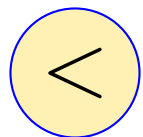
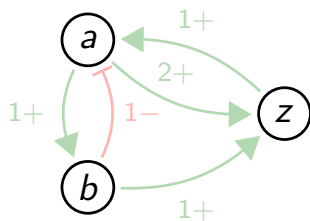
$$z_1 \leftarrow b_1.$$

$$z_0 \leftarrow a_1 \wedge b_0.$$

$$z_0 \leftarrow a_0 \wedge b_0.$$

Discrete Models as Logic Programs

Discrete model:



Expressivity

Logic program:

$$b_1 \leftarrow a_1.$$

$$b_1 \leftarrow a_2.$$

$$b_0 \leftarrow a_0.$$

$$z_1 \leftarrow a_2.$$

$$z_1 \leftarrow b_1.$$

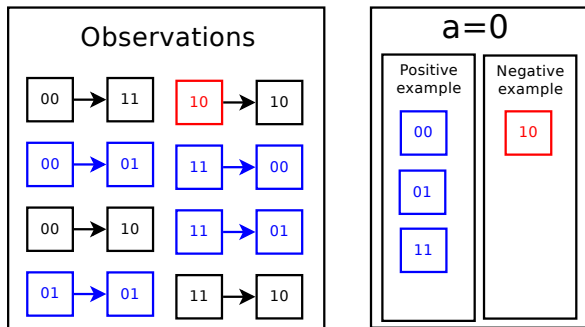
$$z_0 \leftarrow a_1 \wedge b_0.$$

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Learning From Interpretation Transition (LFIT)

Learning Algorithm Intuition: Classification Problem

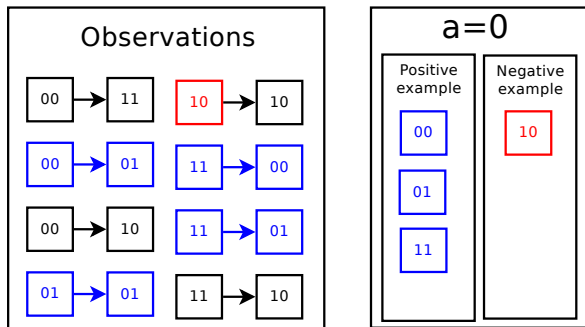
Learn applicable rules: conditions so that a variable **can** take a certain value in next state.



Equivalent to a **classification problem**: What is a typical state where a can take value 0 in the next state ? Here: when a_0 or b_1 is present.

Learning Algorithm Intuition: Classification Problem

Learn applicable rules: conditions so that a variable **can** take a certain value in next state.



Equivalent to a **classification problem**: What is a typical state where a can take value 0 in the next state ? Here: when a_0 or b_1 is present.

That is: $a_0 \leftarrow a_0$. $a_0 \leftarrow b_1$.

Presentation of GULA

GULA = General Usage LFIT Algorithm

Input: a set of transitions ($s_1 \rightarrow s_2$)

Output: a logic program that reproduces the input

Principle: minimal refinements of the rules

Compatible with the synchronous, asynchronous and general semantics
(and any semantics without memory or hard-coded behaviors)

GULA: Initial Logic Program

Suppose:

- a and b have two levels $\{0, 1\}$ and c has three levels $\{0, 1, 2\}$

GULA starts with the most general program:

$$\begin{array}{lll} a_0 \leftarrow . & b_0 \leftarrow . & c_0 \leftarrow . \\ a_1 \leftarrow . & b_1 \leftarrow . & c_1 \leftarrow . \\ & & c_2 \leftarrow . \end{array}$$

With this program, everything is always possible

At each step, some rules can be **refined** by adding atoms to their body

GULA: One Step of Minimal Refinements

Suppose:

- a and b have two levels $\{0, 1\}$ and c has three levels $\{0, 1, 2\}$
- the current program contains the following rules regarding a_1 :

$$a_1 \leftarrow c_2.$$

$$a_1 \leftarrow b_1.$$

- from state $\langle a_1, b_0, c_2 \rangle$, a_1 is never observed in the next states.

However, the first rule allows this; it is then necessary to make **minimal refinements** in order to make this rule inapplicable:

GULA: One Step of Minimal Refinements

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- a and b have two levels $\{0, 1\}$ and c has three levels $\{0, 1, 2\}$
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$$a_1 \leftarrow a_0, c_2.$$

$$a_1 \leftarrow b_1, c_2.$$

$$a_1 \leftarrow c_2, c_0.$$

$$a_1 \leftarrow c_2, c_1.$$

$$a_1 \leftarrow b_1.$$

(No change)

GULA: One Step of Minimal Refinements

Suppose:

- a and b have two levels $\{0, 1\}$ and c has three levels $\{0, 1, 2\}$
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$$a_1 \leftarrow c_2, c_1.$$

GULA: One Step of Minimal Refinements

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$$a_1 \leftarrow b_1, c_2.$$

$$a_1 \leftarrow b_1.$$

(More general)

GULA: One Step of Minimal Refinements

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Example: Synchronous Semantics

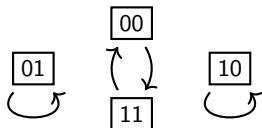


$$f_a = \neg b$$

$$f_b = \neg a$$

b	f_a
0	1
1	0

a	f_b
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1	0



Synchronous

$$// f_a = \neg b$$

$$a_0 \leftarrow b_1$$

$$a_1 \leftarrow b_0$$

$$// f_b := \neg a$$

$$b_0 \leftarrow a_1$$

$$b_1 \leftarrow a_0$$

Example: Asynchronous Semantics

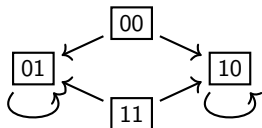


$$f_a = \neg b$$

$$f_b = \neg a$$

b	f_a
0	1
1	0

a	f_b
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1	0



Asynchronous

$$// f_a = \neg b$$

$$a_0 \leftarrow b_1$$

$$a_1 \leftarrow b_0$$

$$// f_b = \neg a$$

$$b_0 \leftarrow a_1$$

$$b_1 \leftarrow a_0$$

// Default rules

$$a_0 \leftarrow a_0$$

$$a_1 \leftarrow a_1$$

$$b_0 \leftarrow b_0$$

$$b_1 \leftarrow b_1$$

GULA: Good Properties

The output of GULA respects some good properties:

- **Consistency**: the program allows no negative examples
- **Realization**: the program covers all positive examples
- **Completeness**: the program covers all the state space
- **Minimality** of the rules (most general conditions)

Results

GULA: an algorithm to learn a biological regulatory network

- From the state graph
- In order to recover the structure of the model
- Applicable to a widespread class of semantics

Limitations:

- Exponential complexity
 - ▶ **PRIDE**: a greedy polynomial version of **GULA**
- What if the data is **incomplete** or **noisy**?
 - ▶ Heuristic to avoid overfitting

Heuristic: Weighted Likelihood/Unlikelihood Rules

- Use the algorithm twice to learn two logic programs:
 - ▶ likelihood rules: what is possible
 - ▶ unlikelihood rules: what is impossible
- Weight each rule by the number of observations it matches

Likelihood rules

$(3, a_0 \leftarrow b_1)$
 $(15, a_1 \leftarrow b_0)$
 $\vdots$

Unlikelihood rules

$(30, a_0 \leftarrow c_1)$
 $(5, a_1 \leftarrow c_0)$
 $\vdots$

Heuristic: Using Weighted Likeliness/Unlikeliness Rules

Explainable predictions:

- Compare weights of applicable likeliness/unlikeliness rules
- Ratio of highest weights $\Rightarrow$ **probability** P
- Rules with highest weights $\Rightarrow$ **explanation** E

$\text{predict} : (atom, state) \mapsto (P, E)$

Likelihood rules

$(3, a_0 \leftarrow b_1)$
 $(15, a_1 \leftarrow b_0)$

Unlikeliness rules

$(30, a_0 \leftarrow c_1)$
 $(5, a_1 \leftarrow c_0)$

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- Ratio of highest weights $\Rightarrow$ **probability** P
- Rules with highest weights $\Rightarrow$ **explanation** E

$\text{predict} : (atom, state) \mapsto (P, E)$

Likelihood rules

$(3, a_0 \leftarrow b_1)$
 $(15, a_1 \leftarrow b_0)$

Unlikeliness rules

$(30, a_0 \leftarrow c_1)$
 $(5, a_1 \leftarrow c_0)$

$\text{predict}(a_1, \langle a_1, b_0, c_0 \rangle) = (0.75, ((15, a_1 \leftarrow b_0), (5, a_1 \leftarrow c_0))) \Rightarrow \text{Likely}$

Heuristic: Using Weighted Likeliness/Unlikeliness Rules

Explainable predictions:

- Compare weights of applicable likeliness/unlikeliness rules
- Ratio of highest weights $\Rightarrow$ **probability** P
- Rules with highest weights $\Rightarrow$ **explanation** E

$\text{predict} : (atom, state) \mapsto (P, E)$

Likelihood rules

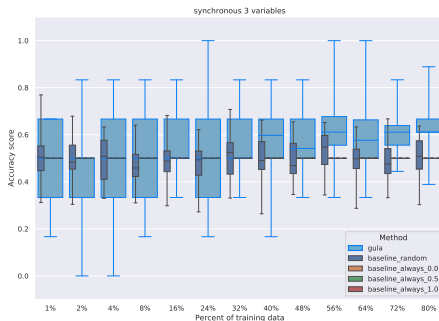
$(3, a_0 \leftarrow b_1)$
 $(15, a_1 \leftarrow b_0)$

Unlikeliness rules

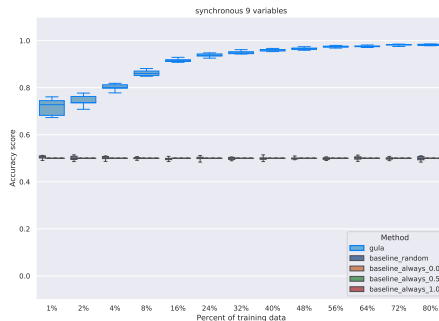
$(30, a_0 \leftarrow c_1)$
 $(5, a_1 \leftarrow c_0)$

$\text{predict}(a_1, \langle a_1, b_0, c_0 \rangle) = (0.75, ((15, a_1 \leftarrow b_0), (5, a_1 \leftarrow c_0))) \Rightarrow \text{Likely}$
 $\text{predict}(a_0, \langle a_1, b_1, c_1 \rangle) = (0.09, ((3, a_0 \leftarrow b_1), (30, a_0 \leftarrow c_1))) \Rightarrow \text{Unlikely}$

Prediction power



3 variables



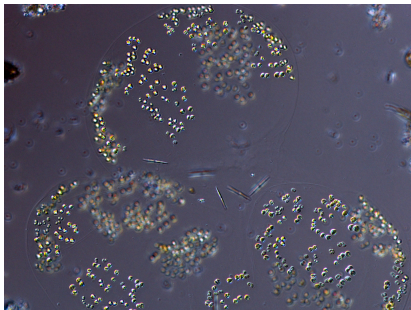
9 variables

Training data = X% of transitions

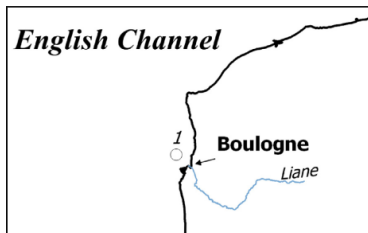
Tested against unseen states (not in the training data)

Application: Dynamics of Marine Phytoplankton

Phytoplankton Blooms



SRN Dataset



<https://www.seanoe.org/data/00397/50832/>

Sampling location	Sampling date	Taxon/Measurment	Value	Sampling depth
001-P-015	1992-05-18	CHLOROA	6.0	Surface (0-1m)
006-P-001	2019-12-02	Chaetoceros	1000.0	Surface (0-1m)
002-P-007	1994-05-25	Pleurosigma	100.0	Surface (0-1m)
002-P-030	2005-10-19	SALI	34.83	Surface (0-1m)
006-P-007	2015-09-28	Guinardia delicatula	11400.0	Surface (0-1m)

Environmental variables (7)

Phytoplankton species (12)

Applying LFIT

Expectations

- Find known **abiotic** influences (of environment on phytoplankton)
- Find new **biotic** influences (of phytoplankton species on others)

Input

- 12 species and 11 environmental variables
- Data interpolation: 1 data point / month

Output

- Run time = 1h (5 runs of **PRIDE**)
- 20'000 rules
- Model accuracy: **0.86**

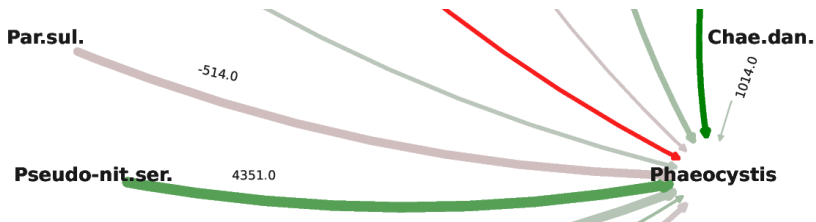
Global Influences

Process: Search and count patterns in rules that characterize an activation/inhibition

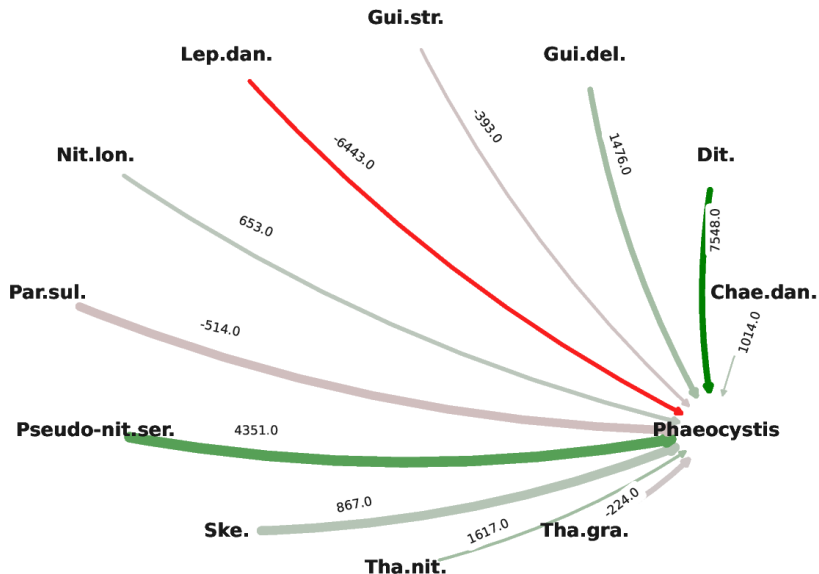
$$a_1 \leftarrow a_0, b_0, c_1.$$

Extracted patterns: $a \xrightarrow{-} a$, $b \xrightarrow{-} a$, $c \xrightarrow{+} a$

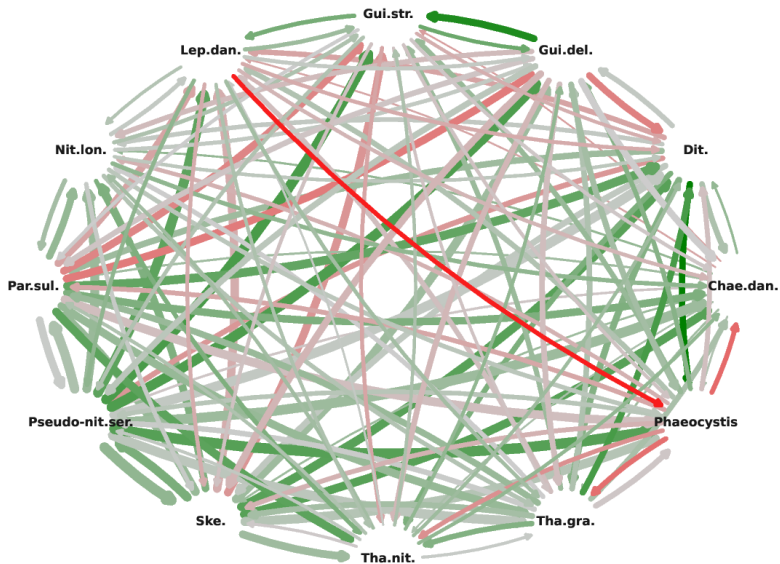
Result: Scores characterizing weight and sign for each pair



Results: Phaeocystis



Results: All Species



Conclusion

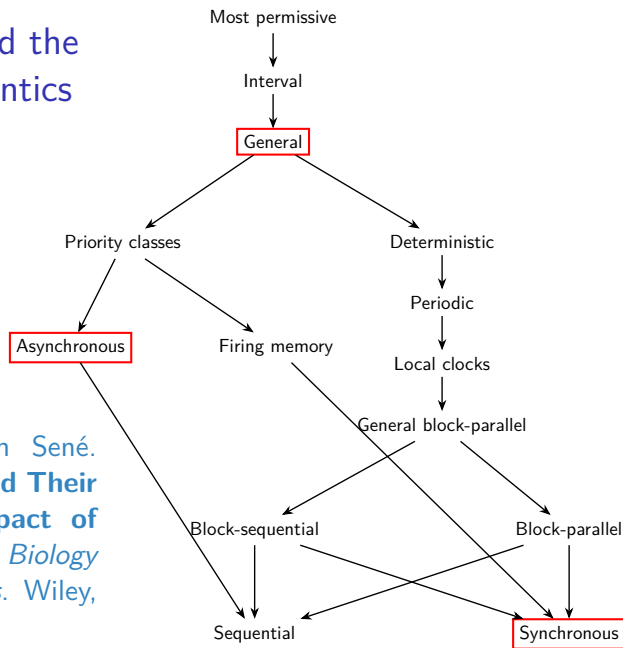
Conclusion

- **Learn** biological regulatory networks with LFIT
- **Heuristics** to tackle real data
 - ▶ Good results with 10% of the transitions
- **Application** to phytoplankton
- You can try LFIT at home:
<https://github.com/Tony-sama/pylfit>

Outlooks:

- Use **GULA** instead of **PRIDE** to get a complete picture
- Improve pre-processing to focus on biotic interactions
- Make LFIT applicable to other semantics

Ongoing: LFIT and the Hierarchy of Semantics



[Loïc Paulevé, Sylvain Sené. **Boolean Networks and Their Dynamics: The Impact of Updates.** In *Systems Biology Modelling and Analysis*. Wiley, 2022.]

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<https://doi.org/10.1007/s10994-021-06105-4>
- **pyLFIT Python library:** <https://github.com/Tony-sama/pylfit>
- **About PRIDE:** Tony Ribeiro, Maxime Folschette, Morgan Magnin and Katsumi Inoue. **Polynomial Algorithm For Learning From Interpretation Transition.** Poster at the *1st International Joint Conference on Learning & Reasoning*. October 2021, Online.
<https://hal.science/hal-03347026v1>
- **Application to phytoplankton:** Madeleine Eyraud, Maxime Folschette, Katsumi Inoue, Sébastien Lefebvre, Cédric Lhoussaine. **Interaction Graphs of Phytoplankton Species Interactions using Logical Learning.** In *The 23st International Conference on Computational Methods in Systems Biology (CMSB25), Lecture Notes in Computer Science*, vol. 15959. September 2025, Lyon, France.